Strong Structural Controllability of Systems on Colored Graphs

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Controllability of systems on graphs

- Consider a simple directed graph $\mathcal{G} = (V, E)$, $V = \{1, ..., n\}$ and $E \subset V \times V$.
- Qualitative class associated with \mathcal{G} :

$$\mathcal{Q}(\mathcal{G}) = \{ A \in \mathbb{C}^{n \times n} \mid \text{for } i \neq j : A_{ji} \neq 0 \Leftrightarrow (i, j) \in E \}.$$

- The diagonal entries of A ∈ Q(G) do not depend on the structure of G, and these are 'free elements' (0 or nonzero).
- Let V_L = {v₁,..., v_m} ⊂ V be the leader set. Associated input matrix B = (e_{v1},..., e_{vm}).
- The leader/follower system defined on graph \mathcal{G} : $\dot{x} = Ax + Bu$, with $A \in \mathcal{Q}(\mathcal{G})$.

Definition: (\mathcal{G}, V_L) is called controllable if (A, B) is controllable for all $A \in \mathcal{Q}(\mathcal{G})$.

Theorem (Monshizadeh, Zhang and Camlibel [1]) Let \mathcal{G} be a simple directed graph and $V_L \subset V$. Then, (\mathcal{G}, V_L) is controllable if and only if V_L is a zero forcing set.



Figure: Simple directed graph \mathcal{G} .

 If v is a black vertex in G with exactly one white out-neighbor u, then we change the color of u to black, and write v → u.

Zero forcing set and controllability of (\mathcal{G}, V_L)

Theorem (Monshizadeh, Zhang and Camlibel [1]) Let \mathcal{G} be a simple directed graph and $V_L \subset V$. Then, (\mathcal{G}, V_L) is controllable if and only if V_L is a zero forcing set.



- If v is a black vertex in G with exactly one white out-neighbor u, then we change the color of u to black, and write v → u.
- The set {1} is a **zero forcing set** for the graph *G*.



Figure: Graph G with $V_L = \{1, 2, 3\}$

- V_L is not a zero forcing set, which implies that (\mathcal{G}, V_L) is not controllable.
- The situation changes drastically if we impose weights to be equal, but still nonzero.



Figure: Graph G with $V_L = \{1, 2, 3\}$

- This graph turns out to be controllable, i.e., (A, B) is controllable for all A associated with this graph. This can be checked by the **PBH** test.
- Note that it becomes very complicated to use the **PBH** test if the graph is of large scale.

Motivation

- Physical: dependence of nonzero parameters
- Particular structure: symmetric, undirected or unweighted



(Picture from "Secure and Efficient Capability-Based Power Management in the Smart Grid")



Figure: Underlying graph with identical edge weights.

Problem formulation

- Colored simple directed graph G(π) = (V, E, π) associated with a simple directed graph G = (V, E) and an edge partition π = {E₁, E₂,..., E_k}, where E_i ∩ E_j = Ø for i ≠ j, and edges have the same color if and only if they are in the same cell.
- Colored qualitative class of G(π) = (V, E, π): a family of matrices associated with G(π):

$$\mathcal{Q}_{\pi}(\mathcal{G}) = \{ A \in \mathcal{Q}(\mathcal{G}) \mid A_{ji} = A_{lk} \ ext{if } (i,j), (k,l) \in E_r ext{ for some } r \}.$$

- $(\mathcal{G}(\pi), V_L)$ is called controllable if (A, B) is controllable for all $A \in \mathcal{Q}_{\pi}(\mathcal{G})$.
- Aim: Establish graph theoretical tests for controllability of $(\mathcal{G}(\pi), V_L)$.









Bipartite graphs and Perfect matching

- A simple directed graph $\mathcal{G} = (V, E)$ is called **bipartite** if there exist two nonempty disjoint vertex sets X and Y such that $X \cup Y = V$ and $(i, j) \in E$ only if $i \in X$ and $j \in Y$, denoted by G = (X, Y, E).
- A set of t edges E ⊂ E is called a t-matching in G, if no two edges in E share a vertex.
- In the special case of |X| = |Y|, such a t-matching is called a perfect matching if t = |X|.



- A simple directed graph G = (V, E) is called **bipartite** if there exist two nonempty disjoint vertex sets X and Y such that X ∪ Y = V and (i, j) ∈ E only if i ∈ X and j ∈ Y, denoted by G = (X, Y, E).
- A set of t edges E ⊂ E is called a t-matching in G, if no two edges in E share a vertex.
- In the special case of |X| = |Y|, such a t-matching is called a perfect matching if t = |X|.
- The pattern class $\mathcal{P}(G)$ of this bipartite graph G is defined as

$$\mathcal{P}(G) = \{ M \in \mathbb{C}^{t \times s} | M_{ji} \neq 0 \Leftrightarrow (x_i, y_j) \in E \}.$$

Colored bipartite graph and colored pattern class

- A colored bipartite graph $G(\pi) = (X, Y, E, \pi)$ is a bipartite graph with a partition of the edge set $\pi = \{E_1, E_1, \dots, E_t\}$ with to each cell assigned a color.
- The pattern class of the colored bipartite graph $G(\pi)$ is defined by

$$\mathcal{P}_{\pi}(G) = \{A \in \mathcal{P}(G) | A_{ji} = A_{lk} \\ \text{if } (x_i, y_j), (x_k, y_l) \in E_r \text{ for some } r\}.$$

Aim: Suppose |X| = |Y|, find a condition to verify whether all A ∈ P_π(G) are nonsingular.
This condition is used to establish the graph-theoretic condition.

Equivalence classes of perfect matchings



Equivalence classes $\mathbb{P} = \{\mathbb{P}_1, \mathbb{P}_2\}$, where $\mathbb{P}_1 = \{p_1, p_3\}$ with $sgn(\mathbb{P}_1) = 0$, and $\mathbb{P}_2 = \{p_2\}$ with $sgn(\mathbb{P}_2) = -1$.

Lemma 1: Let $G(\pi) = (X, Y, E, \pi)$ be a colored bipartite graph and |X| = |Y|. Then, all $A \in \mathcal{P}_{\pi}(G)$ are nonsingular if and only if there exists exactly one equivalence class of perfect matchings having nonzero signature.

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Let $\mathcal{G}(\pi) = (V, E, \pi)$ be a colored simple graph, with each vertex colored either white or black. Let $C \subseteq V(G)$ be set of vertices initially colored black.

- Given X ⊂ C and Y = N_{V\C}(X), assume |X| = |Y|. Note that the subgraph G(X, Y, E_{XY}, π) is a colored bipartite graph. If exactly one equivalence class of perfect matchings has nonzero signature in G(X, Y, E_{XY}, π). Then X forces Y to be black, and we write X ⇒ Y.
- The derived set $D_c(C)$ is the set of black vertices obtained by applying this color change rule until no more changes are possible. If $D_c(C) = V$, C is called a colored zero forcing set.

Example for colored zero forcing set

Theorem 2

Consider a colored simple directed graph $\mathcal{G}(\pi) = (V, E, \pi)$, with leader set $V_L \subset V$. Assume V_L is a colored zero forcing set. Then $(\mathcal{G}(\pi), V_L)$ is controllable.



Note the colored bipartite graph $G(\pi)$ is same as the one in previous example.

 $\begin{array}{l} \mbox{Colored color change rule:} \\ \{1,2,3\} \stackrel{C}{\Rightarrow} \{4,5,6\}. \end{array}$

 $V_L = \{1, 2, 3\}$ is a colored zero forcing set.

 $(\mathcal{G}(\pi), V_L)$ is controllable by Theorem 2.

Main results

Lemma 1 Let $G(\pi) = (X, Y, E, \pi)$ be a colored bipartite graph and |X| = |Y|. Then, all $A \in \mathcal{P}_{\pi}(G)$ are nonsingular if and only if there exists exactly one equivalence class of perfect matchings having nonzero signature.

Theorem 2

Consider a colored simple directed graph $\mathcal{G}(\pi) = (V, E, \pi)$, with leader set $V_L \subset V$. Assume V_L is a colored zero forcing set. Then $(\mathcal{G}(\pi), V_L)$ is controllable.

Corollary 3

Consider a system defined on a colored simple directed graph $\mathcal{G}(\pi) = (V, E, \pi)$ with $|\pi| = |E|$ with leader set $V_L \subset V$. Then (\mathcal{G}, V_L) is controllable if and only if V_L is a colored zero forcing set.

Note: In the special case $|\pi| = |E|$, our condition is equivalent to that in terms of zero forcing set.

- Extension to systems defined on graphs with identical edge weights, which is a other kind of dependencies between the nonzero parameters.
- Establish a graph theoretic test for controllability of system on colored graph in terms of colored zero forcing set.



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Thank you for your attention! The End